

OFF-CENTRES AND RESTRICTED DELAUNAY TESSELLATION:

HIGH-QUALITY MESH GENERATION FOR GENERAL PLANAR, SURFACE AND VOLUMETIC GEOMETRIES

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THE JIGSAW MESHING LIBRARY.

JIGSAW is a new Delaunay-based isotropic mesh generation package designed for general two- and three-dimensional meshing problems [1,2,3]. This new environment seeks to combine a number of useful features:

1. A flexible *geometry-agnostic* formulation based on *restricted* Delaunay tessellations, supporting generalised geometry inputs including piecewise linear complexes, analytic forms and implicit representations.
2. A hybrid *Frontal-Delaunay* refinement strategy – seeking to achieve very high-quality Delaunay-based mesh generation through use of a hybrid, advancing-front type strategy. Using an appropriate set of *off-centre* point-placement rules, this approach aims to combine the best features of classical Delaunay-refinement and advancing-front type techniques, generating high-quality Delaunay meshes that satisfy a set of theoretical bounds and convergence guarantees.
3. A new, publicly-available meshing interface, providing access to a suite of two- and three-dimensional Delaunay-type meshing algorithms via an easy-to-use **MATLAB / OCTAVE** scripting interface.

RESTRICTED DELAUNAY TESSELLATION.

Geometrical features are approximated using a hierarchy of Delaunay sub-complexes [4,5,6]. The bounding Delaunay complex $DEL(X)$ contains the restricted sub-complexes $DEL_{\Gamma}(X)$, $DEL_{\Sigma}(X)$ and $DEL_{\Omega}(X)$, providing discrete approximations to the curve network Γ , the surface patches Σ and the interior volume Ω . $DEL_{\Gamma}(X)$, $DEL_{\Sigma}(X)$ and $DEL_{\Omega}(X)$ contain the 1-, 2- and 3-simplices that best approximate the input curve, surface and volume features, respectively.

Restricted Delaunay techniques exploit the duality between Delaunay tessellations and Voronoi complexes, computing membership for $DEL_{\Gamma}(X)$, $DEL_{\Sigma}(X)$ and $DEL_{\Omega}(X)$ by evaluating intersections between $VOR(X)$ and the input features Γ , Σ and Ω . Such intersections induce a set of circumscribing balls $B(C_i, R_i)$ associated with the faces of $DEL(X)$, providing a set of element-wise refinement points C_i , and a bound on surface discretisation error (Hausdorff distance).

A RESTRICTED DELAUNAY-REFINEMENT ALGORITHM.

Following conventional methodology [4,5,6], a coarse initial Delaunay triangulation is progressively refined until all elements are *good*. At each step, any that violate a set of constraints are identified and the worst offending elements *eliminated*. Elimination is achieved through the insertion of additional *Steiner-vertices* located at the *refinement-points* associated with the elements in question.

```

FUNCTION DELAUNAY-MESH( $\Gamma, \Sigma, \Omega, DEL_{\Gamma}(X), DEL_{\Sigma}(X), DEL_{\Omega}(X)$ )
  WHILE (BAD(E) OR BAD(F) OR BAD(T))
    IF (BAD(E)) THEN
      Refine "worst" restricted edge Ei. Insert "refinement-point"
      for Ei and update  $DEL_{\Gamma}(X)$ ,  $DEL_{\Sigma}(X)$ ,  $DEL_{\Omega}(X)$ .
    ELSE
      IF (BAD(F)) THEN
        Refine "worst" restricted face Fi. Insert "refinement-point"
        for Fi and update  $DEL_{\Gamma}(X)$ ,  $DEL_{\Sigma}(X)$ ,  $DEL_{\Omega}(X)$ . Careful
        not to encroach edges in E.
      ELSE
        IF (BAD(T)) THEN
          Refine "worst" restricted tria Ti. Insert "refinement-point"
          for Ti and update  $DEL_{\Gamma}(X)$ ,  $DEL_{\Sigma}(X)$ ,  $DEL_{\Omega}(X)$ . Careful
          not to encroach edges in E or faces in F.
        END IF
      END WHILE
    END FUNCTION
  
```

Elements are considered to be *bad* based on a combination of geometrical and topological constraints, including radius-edge ratios, surface-discretisation error measures, element size metrics and local topological considerations.

FRONTAL-DELAUNAY METHODS AND "OFF-CENTRES".

Frontal-Delaunay algorithms are a hybridisation of advancing-front and Delaunay-refinement techniques, in which a Delaunay triangulation is used to define the topology of a mesh while new vertices are inserted according to an advancing-front type approach.

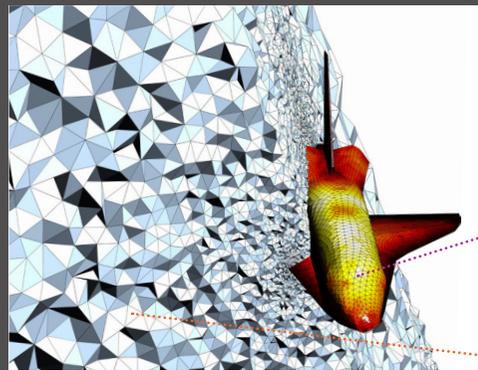
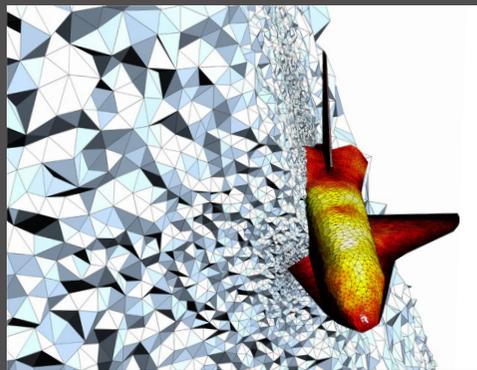
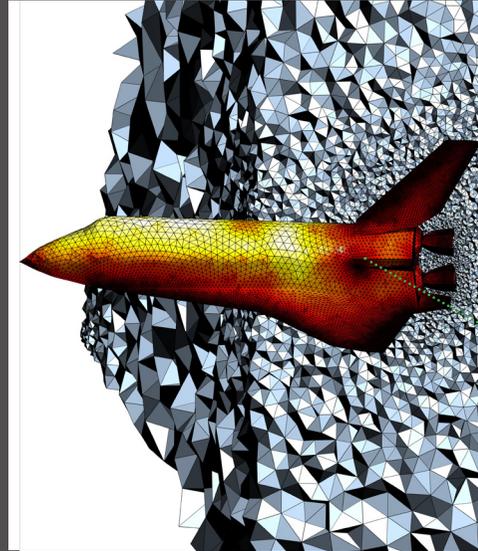
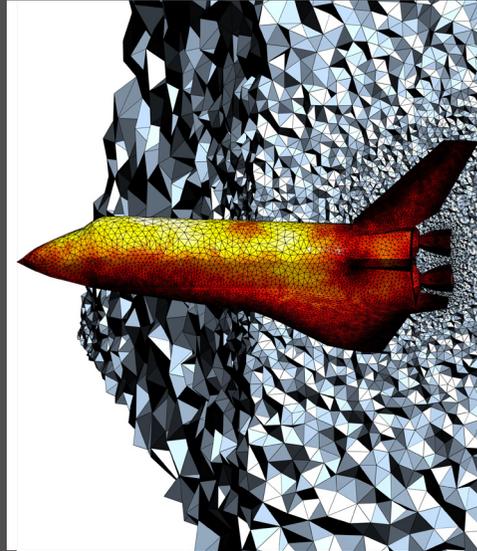
Standard Delaunay-refinement techniques can be transformed to Frontal-Delaunay type algorithms through the selection of appropriate *off-centre* refinement rules [7,1,2,3]:

- 1A. **Edge-refinement:** A locally *size-optimal* point can be formed by computing intersections between a ball of radius $H(X)$ and the curve network Γ . Balls are centred at a frontal vertex. Such a strategy is designed to satisfy local mesh-size constraints.
- 2A. **Face-refinement:** A locally *size-optimal* point can be found by computing intersections between a disk of radius $H(X_m)$ and the surface patches Σ . The disk is centred on a *frontal* edge of the associated face. Such a strategy is designed to satisfy local mesh-size constraints.
- 2B. **Face-refinement:** A locally *shape-optimal* point can be found by computing intersections between a disk of radius ρ_{EMIN} and the surface patches Σ . The disk is centred on a *frontal* edge of the associated face. This construction attempts to satisfy local radius-edge constraints.
- 3A. **Tria-refinement:** A locally *size-optimal* point can be found by computing intersections between a ball of radius $H(X_m)$ and a local Voronoi vector. Such a strategy is designed to satisfy local mesh-size constraints.
- 3B. **Tria-refinement:** A locally *shape-optimal* point can be found by computing intersections between a ball of radius ρ_{VMIN} and a local Voronoi vector. This construction attempts to satisfy local radius-edge constraints.

In each case, it is important to limit off-centre selection to a local *safe-region* – an adjacent sub-region of the Voronoi complex bounded by local element intersections. This constrained point-selection process preserves the convergence guarantees of the underlying Delaunay-refinement algorithm by *falling-back* to standard circumcentre-based refinement in limiting configurations.

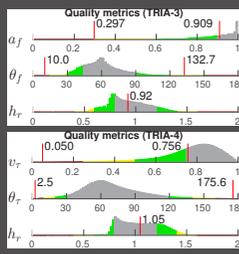
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JIGSAW-DR

Meshes for the SHUTTLE geometry using the Delaunay-refinement algorithm, showing normalised histograms of element volume-length ratio, dihedral-angle and relative edge-length. The mesh contains 217214 surface triangles and 3,644,738 tetrahedrons. Comparing results, it can be seen that the Delaunay-refinement algorithm generates slightly larger meshes with degraded element quality characteristics and mesh-size conformance. Degradation in element volume-length and angle distributions is especially marked in the case of surface triangles, though a similar reduction is also evident in the case of tetrahedral elements. Comparisons of distributions of element relative-length reveal the largest relative differences between algorithms, with the Delaunay-refinement approach displaying significantly weaker conformance to the imposed mesh-size function.



JIGSAW-FD

Meshes for the SHUTTLE geometry using the Frontal-Delaunay algorithm, showing normalised histograms of element volume-length ratio, dihedral-angle and relative edge-length. The mesh contains 191,470 surface triangles and 3,436,473 tetrahedrons. Comparing results, it can be seen that the Frontal-Delaunay algorithm generates slightly smaller meshes with improved element quality characteristics and mesh-size conformance. Improvements in element volume-length and angle distributions is especially marked in the case of surface triangles, though a similar enhancement is also evident in the case of tetrahedral elements. Comparisons of distributions of element relative-length reveal the largest relative differences between algorithms, with the Frontal-Delaunay approach displaying significantly improved conformance to the imposed mesh-size function.

